# Simulation of Conditional Value-at-Risk via Parallelism on Graphic Processing Units

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Simulation of Conditional Value-at-Risk via P

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## **Outline of Topics**

- Motivation
- Sequential procedure of CVaR via nested simulation
- Paralleled nested simulation with screening
- Experimental performance
- Conclusion & Question

#### **Motivations**

- Some simulation procedures take days, weeks or even months to run.
- Traditional supercomputers are inaccessible for most researchers and practitioners.
- New development of high performance computing catches out attention:
  - cloud computing: heterogeneous, massive storage, massive computing
  - parallel computing on GPUs: homogeneous, data light, moderate computing
- Characteristics of simulation in financial engineering: data security, moderate size computing.

# **GPU** Computing

Characteristics of GPU Computing: same program on different data, which makes it attractive for "embarrassingly paralleled" simulation work.

Reported work shows

- 20% 800% faster for sorting (key only)
- 30-50 times faster for random number generating
- almost 100 faster for Monte Carlo simulation of Black Shorel model

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How about its effectiveness on complex simulation such as nested simulation?

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#### **Risk Measurement**

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- $VaR_p = F_V^{-1}(p)$  is the *p*-quantile of  $F_V$ .
- $\operatorname{CVaR}_p = \int_p^1 F_V^{-1}(s) ds.$

#### Nested Simulation

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#### **Nested Simulation**

Goal: provide a framework to estimate the market risk of an investment portfolio

- outer level: generating Z(T) from market scenario model in the real world probability measure.
   underlying stock price S<sub>i</sub>(T), i = 1, · · · , k.
- inner level: estimating V conditionally on Z(T) by Monte Carlo Simulation in risk neutral probability measure. generate the underlying stock prices S(T ≤ t ≤ U)|S<sub>i</sub>(T), i = 1, ..., k

## "Plain" Nested Simulation



- Simulate scenarios Z<sub>1</sub>, · · · , Z<sub>k</sub>. Unknown true value V<sub>i</sub> = E[X|Z = Z<sub>i</sub>].
  Simulate payoffs X<sub>i1</sub>, · · · , X<sub>in₀</sub>
  - given  $Z_i$ . Estimate  $V_i$  by  $(X_{i1} + \cdots + X_{in_0})/n_0$ .

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# Nested Simulation with Screening



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- The comparison of Scenario i and j is taken as a hypothesis test:

$$H_o: V_i > V_j \text{ v.s.} H_1: V_i \leq V_j$$

we say, j is beaten by i if  $H_o$  is true.

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- Screening can be taken as multiple comparisons of each pairs of market scenarios. A market scenario will survive from screening only when it is beaten less than [kp] times of such comparisons.
- Restarting is used to avoid the selection bias inherited from screening.

Embarrassingly Paralleled: Generating Scenarios and 1st Stage Conditional Payoffs

Random Number Generator on GPU:

- long period, multiple streams and sub-streams.
- 32 bits computation only
- combined multiple recursive generator

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Experiments show 50-60 times faster than the sequential procedure.

## Parallel Screening 1: Strategy

Sequential procedure:

• for any scenario  $i = 1, \dots, k$ , compute at least  $\lceil kp \rceil$  variances of difference  $S_{ij}^2 = \sum_{l=1}^{n_0} \frac{(X_{il}-X_{jl})^2}{n_0-1} - \frac{\sum_{l=1}^{n_0} (X_{il}-X_{jl}) \sum_{l=1}^{n_0} (X_{il}-X_{jl})}{n_0(n_0-1)}$ 

• scenario *i* is screened out if it is beaten no less than  $\lceil kp \rceil$  times. thus, a lot of data communication has to be done. To save data communication, we adopt a "divide-and-conquer" strategy.

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  - randomly group the total k market scenario into G groups, each GPU processes one group.
  - apply screening procedure within each group to eliminate some unimportant market scenarios.
  - collect remaining scenarios from each GPU and perform screening again.

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## Parallel Screening 2: Flow Chart



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## Parallel Screening 3: Implementation on GPU



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# Parallel Screening 3: Implementation on GPU



Scheme 1 is about 2 times faster than CPU only, scheme 2 is 30-50 times faster.

# Second Stage Sampling



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# Second Stage Sampling



balance the uneven load among different GPUs and threads

• experiments show 100-400 times faster

## **Experimental Setting**

- hardware: one 3.0G Hz Core 2 Duo CPU, four GeForce 9800 GPUs, 6 GB memory.
- software: gcc 4.2.4 on Ubuntu 64 bits version.
- testing portfolios: shorting one put option and a portfolio with 8 options.
- number of replications: 100 and more.

# Speedup of Sampling



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## Speedup of Computing the CI



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# Speedup in Running Time



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- Parallelism on GPU is very similar as that on CPUs and distributed computing, data communication and synchronization are the bottlenecks of scalability.
- Research opportunities exist in providing standard scientific computing package like BLAST on GPU computing.



# Multi-GPU and OpenMP

- OpenMP
  - shared memory(fast cached)
  - high performance yet limited scalability
  - branching adaptability
- CUDA on GPU
  - many cores
  - shared memory (cached and un-cached)
  - fast on-board memory (register, local, share, global, const ,texture)
  - hardware thread management (zero overhead)
- 4 GPUs allowed on one computer. Affordable personal supercomputer.

# CUDA Program Model



16 successive threads is called half-wrap, processed simultaneously on one SP.

- no diversity
- access successive memory in one command

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The host issues a succession of kernel invocations to the device. Each kernel is executed as a batch of threads organized as a grid of thread blocks

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# Memory Model



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