

Dynamic, Risk-Based Aviation Security Screening Policy Performance Analysis Using Simulation

Sheldon H. Jacobson
Department of Computer Science
University of Illinois at Urbana-Champaign

Joint work with Adrian J. Lee (CITERI)
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Ni Hao
Ni hui shuo yingyu ma?

I hope so!

Research Motivation

- Investigate passenger / baggage screening operations
- Effectively utilize security resources
- Maximize security system effectiveness
- **Design security systems that work!**



Overview

- Introduction
- Historical Background
- Steady State Assignment Policy
 - Maximize passenger throughput
- Transient Assignment Policy
 - Maximize security *and* passenger throughput
- Selective Screening Systems
- Simulation Results
- Observations and Conclusions

Introduction

- Aviation Security: A New Era
 - September 11, 2001
 - Terrorist attacks on World Trade Center Twin Towers and Pentagon
 - Ongoing Events
 - August 10, 2006: Plot to destroy ten U.S. bound transatlantic flights
 - December 25, 2009: Christmas day bomber
 - May 7, 2012: Al-Qaeda bomber
- Transportation Security Administration (TSA)
- Changes to aviation security systems
 - Reinforced cockpit doors
 - Expanded federal air marshal program
 - Behavior detection techniques (SPOT)
 - 100% checked baggage screening
 - 3-1-1 liquid and gel policy
 - Advanced Imaging Technologies



www.tsa.gov

Passenger Screening Techniques

Uniform screening

- Rationale: All passengers could pose a threat
- Passenger risk perceived equally
- Used from 1970's to 1998

Selective screening

- Rationale: Majority of passengers pose no threat
- Select passengers perceived as higher risk
- Targets expensive, specialized resources at high-risk passengers
- Used from 1998 to 2001
- TSA Pre✓ is the new trusted traveler program

Passenger Prescreening Programs

- Computer-Assisted Passenger Prescreening System (CAPPS)
 - **Selectees** - those not cleared by CAPPS
 - **Nonselectees** - those cleared by CAPPS
 - CAPPS II (2003)
 - Lacked proper analysis and tests during development
 - Dismantled due to privacy concerns
- Secure Flight (2004)



- Registered Traveler (RT) programs www.tsa.gov
 - Expedites screening process for RT members (Global Entry, Nexus)
 - TSA Pre✓

Baggage and Cargo Screening

- Commission on Aviation Safety and Security, July 1996
 - Explosive detection systems (EDSs)
 - Automated passenger prescreening (i.e., CAPPS)
 - Positive passenger-baggage matching (PPBM)
- Aviation and Transportation Security Act (ATSA)
 - 100% screening of checked baggage by December 31, 2002
- 9/11 Commission Act of 2007 – Recommendations
 - Required “screening” of
 - 50% of cargo on passenger aircraft by February 2009
 - 100% of cargo by August 2010
 - Security Programs
 - Explosives detection canine teams
 - Transportation Security Inspectors (TSIs) for cargo



www.gesecurity.com



www.tfsrc.gov

Current Passenger Screening Programs

- Checkpoint Evolution
 - People
 - Travel Document Checker (TDC)
 - Visible Intermodal Prevention and Response (VIPR)
 - Screening Passengers by Observation Technique (SPOT)
 - Chat downs
 - Process
 - Diamond Self-Select program (3 groups)
 - TSA Pre✓
 - Technology
 - Advanced Imaging Technology (AIT)
 - Trace Devices
 - Bottle Liquid Scanners (BLS)



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Definitions

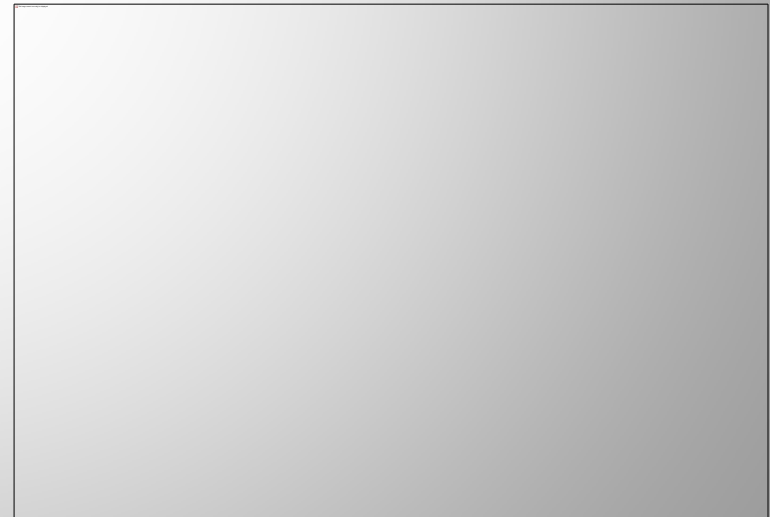
- A **threat** is a passenger or item that may be involved in the attack of an aircraft or within the airport terminal
- A **device** is a technology or procedure used to detect a threat
- The device **capacity** is an upper bound on the number of passengers or bags that a device can screen
- A **security class** is a subset of devices and procedures through which a passenger may be screened
- **Multi-level screening** is an aviation security system in which there exists several security classes to screen passengers
- An **assessed threat value** quantifies the passenger's perceived risk level through an automated prescreening system

Device Alarm Responses (EPIC)

- **True Alarm** (Effective)
 - An alarm occurs for a passenger/bag containing a threat item
 - Correctly identifies a potential terrorist attack
- **False Clear** (Perilous)
 - No alarm occurs for a passenger/bag containing a threat item
 - Incorrectly allows a potential terrorist to enter the airport terminal
- **False Alarm** (Inefficient)
 - An alarm occurs for a passenger/bag containing no threat items
 - Requires additional screening, cost, time
- **True Clear** (Convenient)
 - No alarm occurs for a passenger/bag containing no threat items
 - Correctly clears nonthreatening passenger

Designing Effective Screening Systems

- Challenges
 - Budget limitations, time consuming, questionable effectiveness
- Improving security screening systems
 - Screen high-risk subjects with expensive, low throughput devices
 - Design **layered** approach to screening passengers, baggage and cargo
 - Assign passengers, baggage, cargo based on **perceived risk**
- Modeling Approach
 - **Real-time dynamic model**
 - Maximize security, subject to device constraints

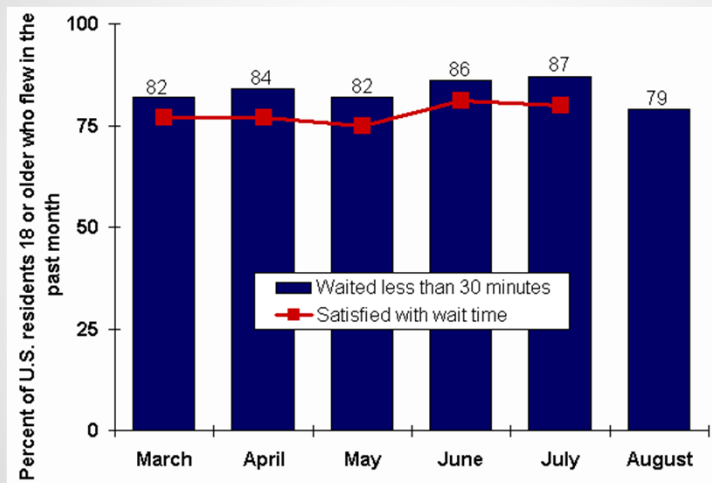


www.tsa.gov

Motivation

Objectives

- Maximize security (overall true alarm rate)
- Minimize expected time passenger spends in security system



Bureau of Transportation Statistics, 2002, www.bts.gov

- Queueing aspects of passenger screening process
- Continual arrival of passengers at security checkpoint ($N \rightarrow +\infty$)
- Multi-level vs. selective screening systems

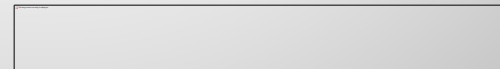
Setting and Notation

- Queue capacity c_m for security class $m = 1, 2, \dots, M$
- Assessed threat value α_i of passenger i
 - Quantifies perceived risk resulting from prescreening
- Conditional probability of security class m detecting threat, L_m (i.e., device true alarm rate)

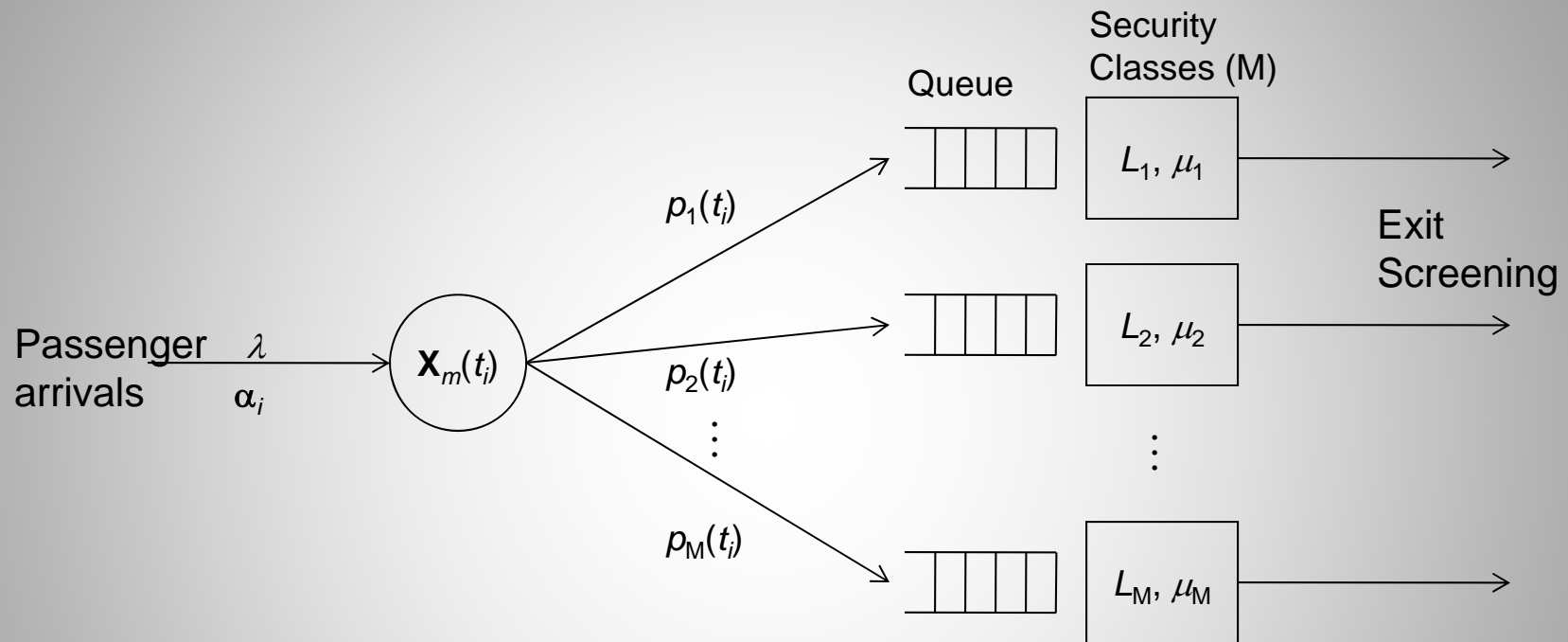
Setting and Notation

- Passenger arrivals
 - Independent
 - Poisson process with rate $\lambda > 0$
 - t_i = time passenger i arrives at security checkpoint ($t_i \rightarrow \infty$ as $i \rightarrow \infty$)
- Security class service times
 - Exponential, with rates $\mu_1 > \mu_2 > \dots \mu_M > 0$
 - Stability $\lambda < \sum_{m=1}^M \mu_m$
- Passenger assignments
 - Probability passenger i assigned to security class m

$$p_m(t_i) \equiv P(\mathbf{X}_m(t_i) = 1)$$



Multi-level Security Class System



- c_m : Queue capacity for security class $m = 1, 2, \dots, M$
- Security classes operate independently

Steady State Assignment Policy

- Passenger assignments made independently
- Predetermined set of fixed security class threshold values
- Each security class has *infinite capacity* ($c_m = +\infty, m = 1, 2, \dots, M$)
- Discrete random variables

– $N^a(t)$ Number of passengers that arrive for screening **by** time t

– $N_m^a(t)$ Number of passengers assigned to class m **by** time t

$$N^a(t) = \sum_{m=1}^M N_m^a(t)$$

– $N_m^d(t)$ Number of passengers screened in class m **by** time t

– $S_m(t)$ Number of passengers in security class m **at** time t

$$S_m(t) = N_m^a(t) - N_m^d(t)$$

One-Step Analysis

- $\{\mathbf{N}_m^a(t), t \geq 0\}$, $m = 1, 2, \dots, M$ independent Poisson processes, with rate λp_m , where $p_m \equiv P(X_m(t_i) = 1)$
- Probability transition rates (Chapman Kolmogorov Equations)

$$\frac{d}{dt} P(S_m(t) = 0) = -\lambda p_m P(S_m(t) = 0) + \mu_m P(S_m(t) = 1)$$

$$\begin{aligned} \frac{d}{dt} P(S_m(t) = s_m) = & -\lambda p_m (P(S_m(t) = s_m) - P(S_m(t) = s_m - 1)) \\ & + \mu_m (P(S_m(t) = s_m + 1) - P(S_m(t) = s_m)), \quad s_m \geq 1 \end{aligned}$$

- Steady-state probabilities

$$P_s^m \equiv \lim_{t \rightarrow \infty} P(S_m(t) = s) = P(S_m = s)$$

- Geometric distribution

$$P_s^m = \left(1 - \frac{\lambda p_m}{\mu_m}\right) \left(\frac{\lambda p_m}{\mu_m}\right)^s, \quad \lambda p_m < \mu_m$$

Expectation and Variance

- Standard Queueing Results for Security Class $m = 1, 2, \dots, M$

- S_m = Steady-state number of passengers in security class m

$$E[S_m] = \frac{\lambda p_m}{\mu_m - \lambda p_m}, \quad Var(S_m) = \frac{\mu_m \lambda p_m}{(\mu_m - \lambda p_m)^2} \quad \lambda_m < \mu_m$$

- W_m = Steady-state amount of time a passenger spends in class m

$$E[W_m] = \frac{1}{\mu_m - \lambda p_m}, \quad Var(W_m) = \frac{\mu_m}{\lambda p_m (\mu_m - \lambda p_m)^2} \quad \lambda_m < \mu_m, \quad p_m \neq 0$$

➤ Security system

$$E[S] = \sum_{m=1}^M E[S_m], \quad Var(S) = \sum_{m=1}^M Var(S_m)$$

$$E[W] = \sum_{m=1}^M p_m E[W_m], \quad Var(W) = \sum_{m=1}^M p_m^2 Var(W_m)$$

- Service rate mean, variance:

$$\bar{\mu} = \frac{1}{M} \sum_{m=1}^M \mu_m, \quad \sigma_\mu^2 = \frac{1}{M} \sum_{m=1}^M (\mu_m - \bar{\mu})^2$$

Steady State Solutions

- Equalize $E[S_m]$ across $m = 1, 2, \dots, M$

 \Rightarrow

$$p_m = \frac{\mu_m}{M\bar{\mu}}$$

$$E[S] = \frac{M\lambda}{M\bar{\mu} - \lambda} \quad \text{Var}(S) = \frac{M^2\bar{\mu}\lambda}{(M\bar{\mu} - \lambda)^2}$$

$$E[W] = \frac{M}{M\bar{\mu} - \lambda} \quad \text{Var}(W) = \frac{M^2\bar{\mu}}{\lambda(M\bar{\mu} - \lambda)^2}$$

- Equalize $E[W_m]$ across $m = 1, 2, \dots, M$

 \Rightarrow

$$p_m = \frac{1}{\lambda}(\mu_m - \bar{\mu}) + \frac{1}{M}$$

$$E[S] = \frac{M\lambda}{M\bar{\mu} - \lambda} \quad \text{Var}(S) = \left(\frac{M}{M\bar{\mu} - \lambda} \right)^2 (M\sigma_\mu^2 + \lambda\bar{\mu})$$

$$E[W] = \frac{M}{M\bar{\mu} - \lambda} \quad \text{Var}(W) = \left(\frac{M}{\lambda(M\bar{\mu} - \lambda)} \right)^2 (M\sigma_\mu^2 + \lambda\bar{\mu})$$

Static Passenger Queueing Problem (SPQP)

- Minimize expected passenger security sojourn time, $E[W]$

$$\begin{aligned} \text{minimize} \quad & \sum_{m=1}^M \frac{p_m}{\mu_m - \lambda p_m} \\ \text{subject to} \quad & 0 \leq p_m \leq 1 \quad m = 1, 2, \dots, M \\ & p_m < \mu_m / \lambda \quad m = 1, 2, \dots, M \\ & \sum_{m=1}^M p_m = 1 \end{aligned}$$

- Solve nonlinear program (NLP) for p_1, p_2, \dots, p_M
- Second inequality constraint replaced with $p_m + \varepsilon_m \leq \mu_m / \lambda$
 - Take limit as $\varepsilon_m \rightarrow 0$

Example: Two-Class Security System

- $p_1 = P(X_1(t_i)=1)$, $p_2 = P(X_2(t_i)=1) = 1-p_1$, with $\mu_1 + \mu_2 > \lambda$, $\mu_1 > \mu_2$
- Solution to SPQP

$$p_1^* = \frac{\mu_1(\lambda - 2\mu_2)}{\lambda(\mu_1 - \mu_2)} + \sqrt{\left(\frac{\mu_1(\lambda - 2\mu_2)}{\lambda(\mu_1 - \mu_2)}\right)^2 + \frac{\mu_1\mu_2}{\lambda^2} - \frac{\mu_1(\lambda - 2\mu_2)}{\lambda(\mu_1 - \mu_2)}}$$

$$p_2^* = \frac{\mu_2(2\mu_1 - \lambda)}{\lambda(\mu_1 - \mu_2)} - \sqrt{\left(\frac{\mu_1(\lambda - 2\mu_2)}{\lambda(\mu_1 - \mu_2)}\right)^2 + \frac{\mu_1\mu_2}{\lambda^2} - \frac{\mu_1(\lambda - 2\mu_2)}{\lambda(\mu_1 - \mu_2)}}$$

- If $\mu_1 = \mu_2$, then $p_1^* = p_2^* = 1/2$

Transient Assignment Policy

- Objectives
 - **Maximize** security (true alarm rate)
 - **Minimize** expected passenger security sojourn time
- Queueing analysis
 - Observe the screening process at each passenger arrival time
 - Interarrival times: exponential with rate $\lambda > 0$
 - $\delta_i = t_i - t_{i-1}$
 - **Finite** security class capacities, c_m
- Weighted cost function
 - Minimize cost to create balance between objectives

Transient Analysis

- Number of passengers screened in class m during $(t_i, t_{i+1}]$

$$N_m^s(t_i, t_{i+1}) = N_m^d(t_{i+1}) - N_m^d(t_i)$$

- Independent of passenger arrival time, t_i
- Dependent on number of passengers in the system at time t_i , $\{S_m(t_i)\}$

- Conditional probability for the number of passengers screened

$$P(N_m^s(t_i, t_{i+1}) = n_m^s | S_m(t_i) = s_m) = \begin{cases} e^{-\mu_m \delta_{i+1}} (\mu_m \delta_{i+1})^{n_m^s} / n_m^s! & \text{if } n_m^s < s_m \\ 1 - \sum_{n_m^s=0}^{s_m-1} e^{-\mu_m \delta_{i+1}} (\mu_m \delta_{i+1})^{n_m^s} / n_m^s! & \text{if } n_m^s = s_m \\ 1 & \text{if } s_m = 0 \end{cases}$$

Markov Chain

- Model as discrete-time, inhomogeneous Markov chain

$$P_m^{k,j}(t_i) = \begin{cases} (1-p_m(t_i))P(N_m^s(t_i, t_{i+1})=k|S_m(t_i)=k) & \text{for } k=0,1,\dots,c_m, \quad j=0 \\ + p_m(t_i)P(N_m^s(t_i, t_{i+1})=k+1|S_m(t_i)=k) & \\ p_m(t_i)P(N_m^s(t_i, t_{i+1})=0|S_m(t_i)=k) & \text{for } k=0,1,\dots,c_m-1, \quad j=k+1 \\ (1-p_m(t_i))P(N_m^s(t_i, t_{i+1})=k-j|S_m(t_i)=k) & \text{for } k=1,2,\dots,c_m, \quad 1 \leq j \leq k \\ + p_m(t_i)P(N_m^s(t_i, t_{i+1})=k-j+1|S_m(t_i)=k) & \\ P(N_m^s(t_i, t_{i+1})=0|S_m(t_i)=k) & \text{for } k=j=c_m \\ 0 & \text{otherwise} \end{cases}$$

- Boundary condition

$$P(S_m(t_1)=s_m) = \begin{cases} 1 & \text{if } s_m = 0 \\ 0 & \text{otherwise} \end{cases}$$

- States $s_m = 0,1,\dots,c_m$ positive recurrent and aperiodic

Closed-Form Recursions

- Expected number of passengers in security class m

$$E[S_m(t_{i+1})] = E[S_m(t_i)] + p_m(t_i) \left(1 - P(N_m^s(t_i, t_{i+1}) = 0, S_m(t_i) = c_m)\right) - E[N_m^s(t_i, t_{i+1})]$$

– Boundary condition, $E[S_m(t_1)] = 0$

- Expected amount of time passenger $i+1$ spends in security system if assigned to security class m

$$E[W_m(t_{i+1})] = E[W_m(t_i)] + \frac{p_m(t_i)}{\mu_m} \left(1 - P(N_m^s(t_i, t_{i+1}) = 0, S_m(t_i) = c_m)\right) - \frac{1}{\mu_m} E[N_m^s(t_i, t_{i+1})]$$

– Boundary condition, $E[W_m(t_1)] = 1/\mu_m$

- Security class threshold values

$$b_m(t_i) = F_\alpha^{-1} \left(\sum_{j=1}^m p_j(t_i) \right)$$

Cost Function Components

- False Clears

$$C^Z(t_i) \equiv \left(1 - \sum_{m=1}^M \frac{L_m - L_1}{L_M - L_1} p_m(t_i) \right)^2$$

- Passenger Sojourn Times

$$C^W(t_i) \equiv \left(\frac{\sum_{m=1}^M p_m(t_i) E[W_m(t_i)] - \omega^*}{\max_{m=1,2,\dots,M} \{E[W_m(t_i)]\} - \omega^*} \right)^2$$

- Optimal, steady-state expected amount of time a passenger spends in the security system, $\omega^* = \sum_{m=1,2,\dots,M} p_m^* \omega_m^*$

- Optimal assignment probability error, $p_m(t_i) - p_m^*$

$$C^P(t_i) \equiv \frac{1}{M-1} \sum_{m=1}^{M-1} \left(1 - \frac{p_m(t_i)}{p^*} \right)^2 \quad \text{where } p^* = \max\{p_m^*\}$$

Dynamic Passenger Queueing Problem

- Total Weighted Cost Function

- $0 \leq \eta_1 \leq 1, 0 \leq \eta_2 \leq 1$

- minimize $C(t_i) = (1 - \eta_1)C^Z(t_i) + \eta_1((1 - \eta_2)C^W(t_i) + \eta_2C^P(t_i))$

- subject to $0 \leq p_m(t_i) \leq 1, \quad m = 1, 2, \dots, M$

- $$\sum_{m=1}^M p_m(t_i) = 1$$

- Solve nonlinear program for $p_1(t_i), p_2(t_i), \dots, p_M(t_i)$

Simulation Model

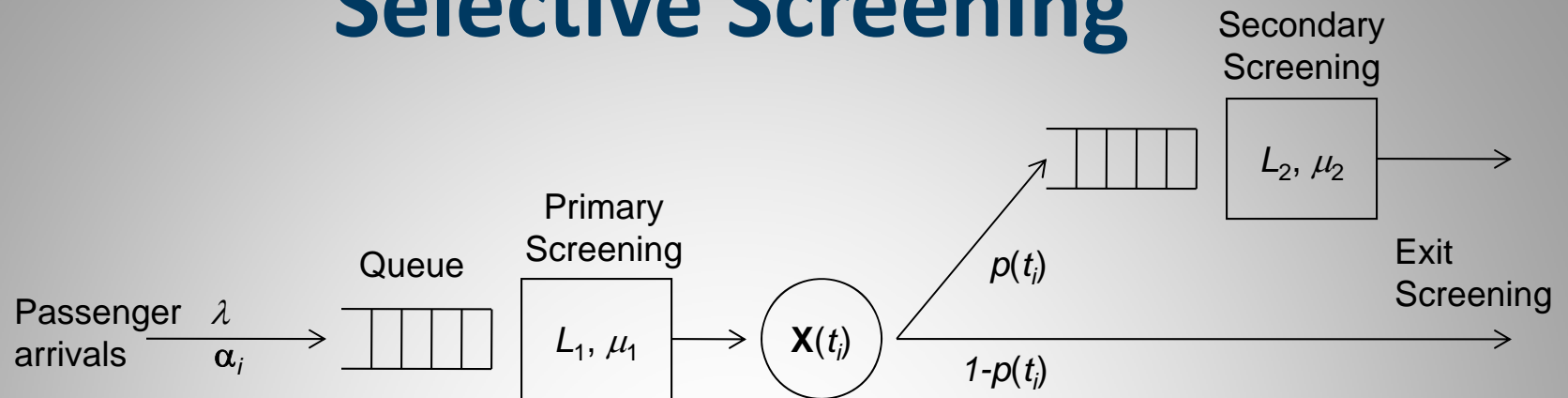
Used to compare security screening policies and conduct sensitivity analysis on the objective function parameters.

Threshold values (used to assigned passengers to classes) are updated each time a passenger arrives, by solving the NLP for the SPQP.

$$b_m(t_i) = F_\alpha^{-1} \left(\sum_{j=1}^m p_j(t_i) \right)$$

Independently seeded runs are used to estimate the mean and the variance of the number of threat items detected and of the time spent within the screening process.

Selective Screening



- Optimal assignment probability for secondary screening

- Nonselectees

$$p^*(t_i) = \begin{cases} \frac{1-\eta_1}{1-p_s} & \text{if } S_2(t_i) < c_2 \\ 0 & \text{otherwise} \end{cases}$$

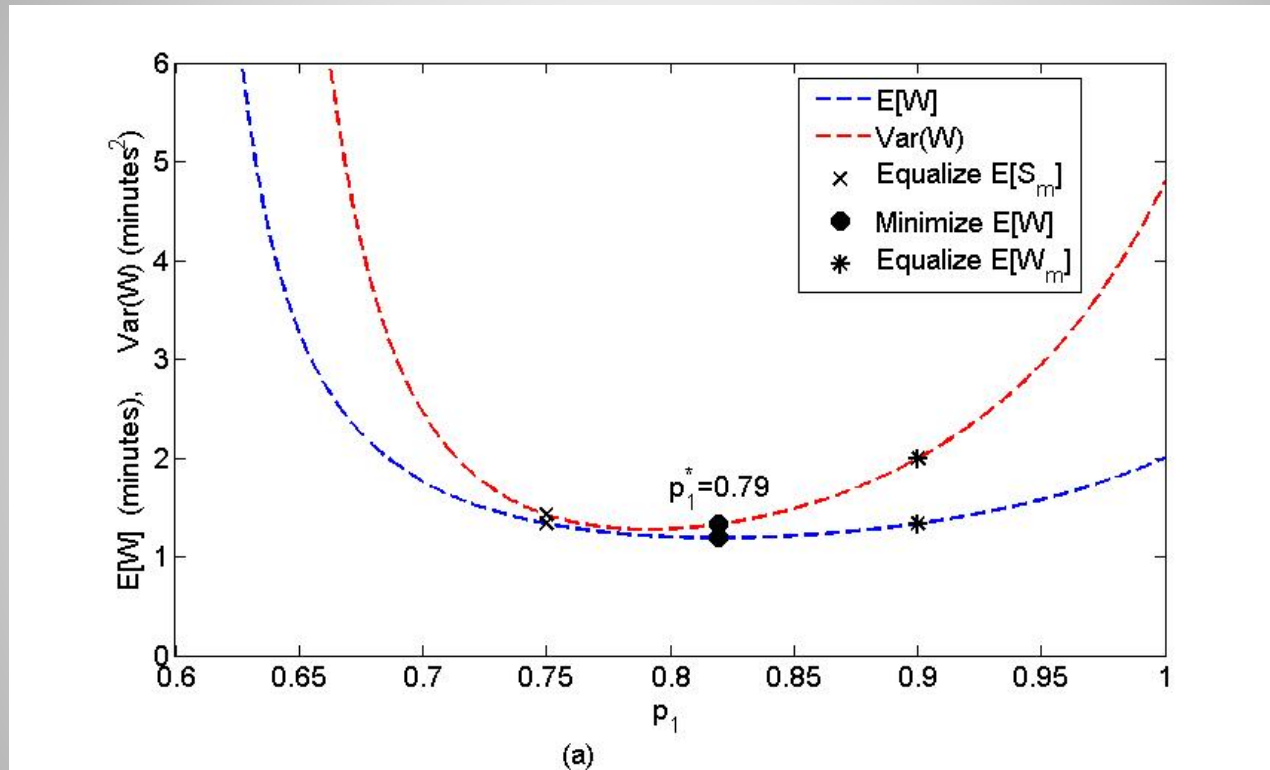
- Selectees, $p^*(t_i) = 1$

- p_s = fraction of passengers designated as selectees

Simulation Results

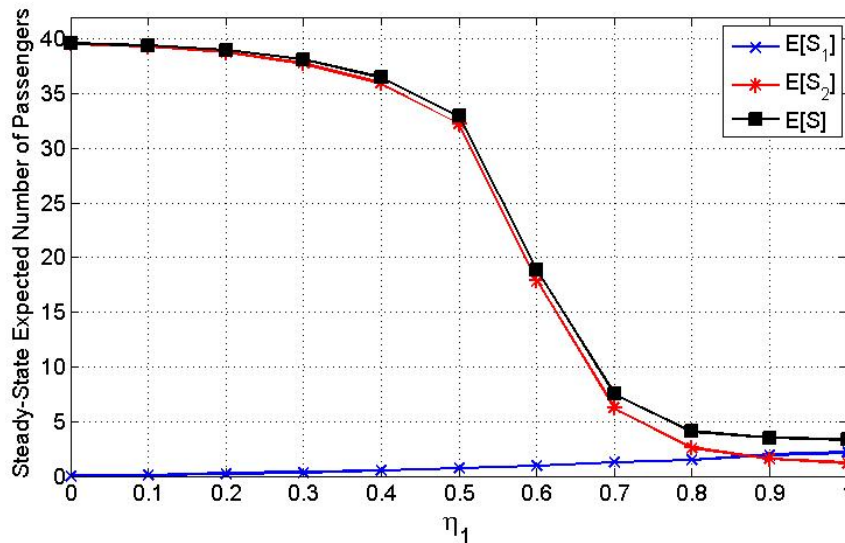
- $M = 2$ class security system, with $N = 1000$ passengers
- $F_{\alpha}(\alpha)$ truncated exponential distribution (over $[0,1]$)
 - $\theta \approx E[\alpha]$ for values $\theta < 0.1$
- Passengers arrive as a Poisson process
 - $\lambda = 2.5$ passengers/minute
- Exponential service times
 - $\mu_1 = 3$ passengers/minute, $\mu_2 = 1$ passengers/minute
- Security levels
 - $L_1 = 0.75$, $L_2 = 0.9$ (class 2 more secure than class 1)
- Capacities
 - $c_1 = 60$, $c_2 = 40$

Static Analysis for M = 2 Class Security System



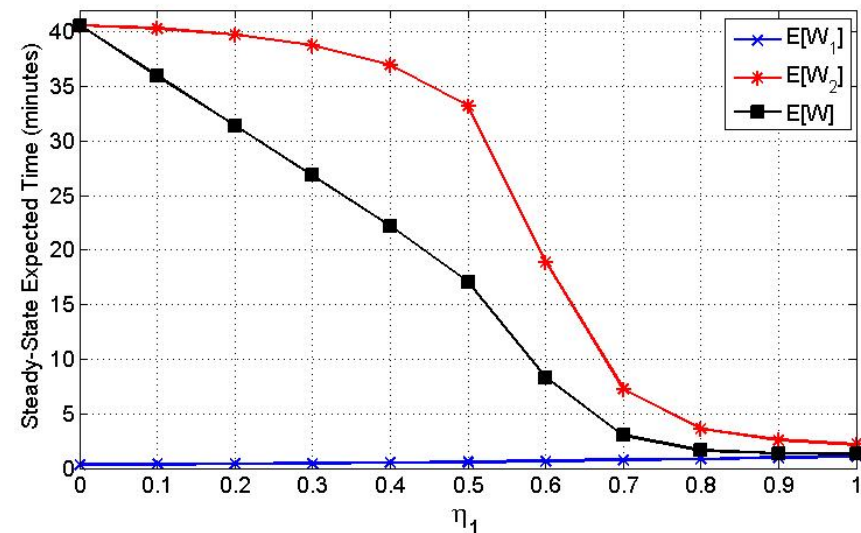
- Minimizing $E[W]$ does not simultaneously minimize $Var(W)$

Dynamic Analysis

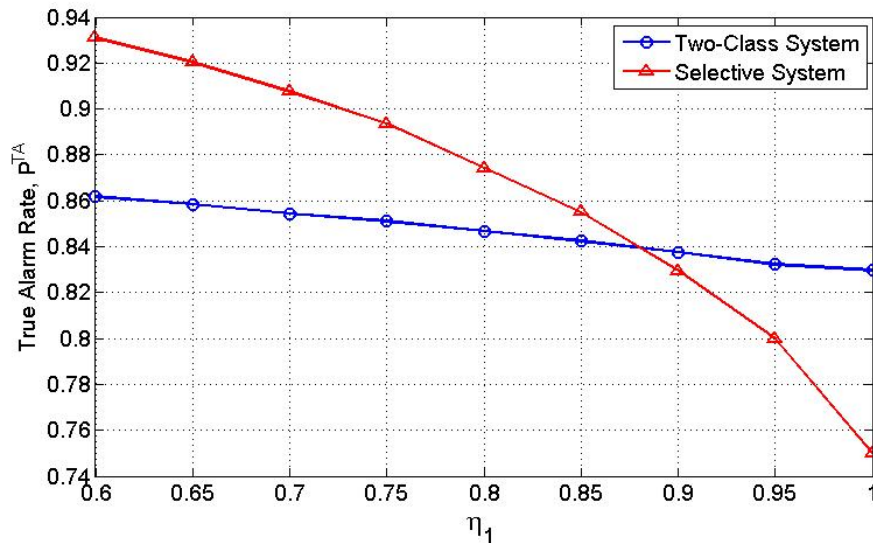


- Effect of cost weight η_1
- Expected number of passengers equalized in each security class when $\eta_1 = 0.9$

- Overall security (true alarms) maximized when $\eta_1 = 0$
- Expected amount of time in security system minimized when $\eta_1 = 1.0$

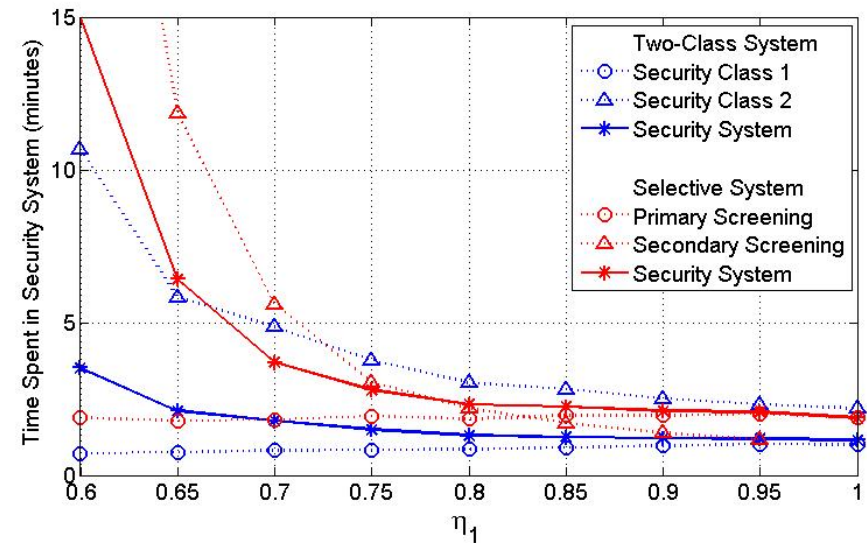


Security System Comparison



- System Security (true alarm rate)
- Selective system designed for maximizing security
- When $\eta_1 = 1$, selective system security is lower since no passengers undergo secondary screening

- Mean time a passenger spends in the security system (minutes)
- Two-class system designed for maximizing passenger throughput



Simulation Contribution

- Simulation is needed to estimate passenger security system sojourn time due to the dynamically evolving security threshold values within the risk-based screening policy.
- Simulation results demonstrate that a **multi-level structure** is designed to expedite screening, while a **selective screening** system increases the probability of detecting threat items.
- Simulation can be used to compare the performance of various alternative security checkpoint designs to analyze the effect on true/false alarm rates and screening times.

Summary

- Aviation Security Application
 - Systematic analysis of the passenger screening process
 - Optimal design of sequential passenger assignment policies
 - Responsively adapt to changing threat environments
- Future Extensions
 - Non-exponential interarrival and service time distributions
 - Explore alternative security system structures
 - Investigate dependency among security classes
 - Incorporate cost associated with resolving alarms (true vs. false)

How to get through airport security without a problem



Thank You (Xie Xie)

Sheldon H. Jacobson
Simulation & Optimization Laboratory
Department of Computer Science

shj@illinois.edu

<https://netfiles.uiuc.edu/shj/www/shj.html>

